

Secure Channels

Summer Term 2018

Problem Set 7

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Agenda

- Auditing Cryptographic Protocols
- Auditing Cryptographic Implementations
- Own Protocols

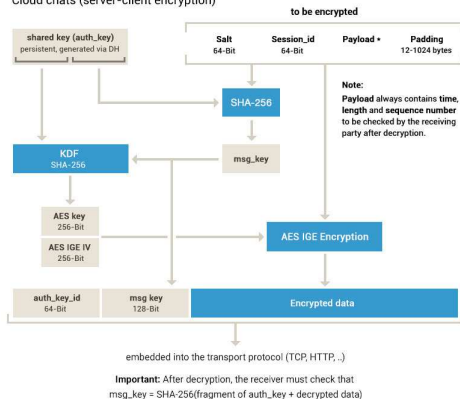
Section 1

Auditing Cryptographic Protocols

Auditing Crypto: MTProto 2

MTProto 2.0, part I

Cloud chats (server-client encryption)



<https://core.telegram.org/mtproto>

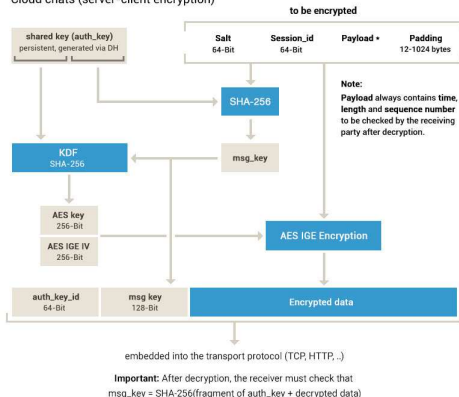
Auditing Crypto: MTPROTO 2

Issues:

- Usage of single-keyed SHA-256 as MAC
- Weak KDF
- KDF is not collision-resistant: 4 key bytes of Auth-key unused
- Key ID is exchangeable
- ...

MTPROTO 2.0, part I

Cloud chats (server-client encryption)



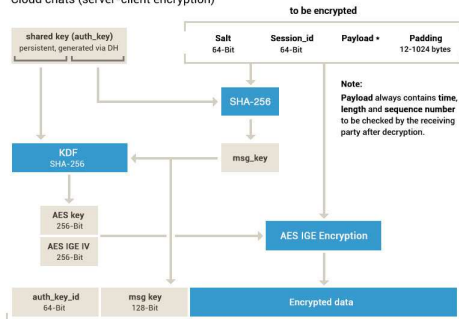
Auditing Crypto: MTPROTO 2

Usage of Weak Cryptographic Schemes

- Single-keyed SHA-256 is not an SUF-CMA-secure MAC
- $F(K || M)$ is vulnerable to extension attacks

MTPROTO 2.0, part I

Cloud chats (server-client encryption)



embedded into the transport protocol (TCP, HTTP, ..)

Important: After decryption, the receiver must check that $msg_key = SHA-256(\text{fragment of auth_key} + \text{decrypted data})$

```
1 | msg_key_large = sha256(self.auth_key[88: 88 + 32] + data + padding).digest()
```

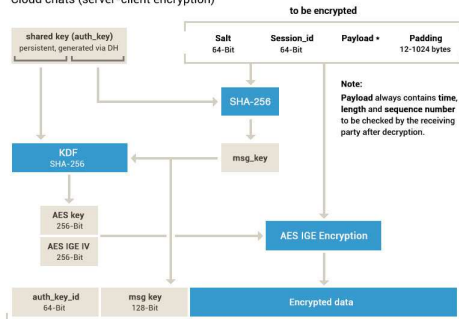
Auditing Crypto: MTPROTO 2

Usage of Weak Cryptographic Schemes

- 2x SHA-256 is not a secure KDF
- Does not thwart offline adversaries

MTPROTO 2.0, part I

Cloud chats (server-client encryption)



embedded into the transport protocol (TCP, HTTP, ..)

Important: After decryption, the receiver must check that
 $msg_key = SHA-256(\text{fragment of auth_key} + \text{decrypted data})$

```
1 sha256_a = sha256(msg_key + auth_key[x: x + 36]).digest()
2 sha256_b = sha256(auth_key[x + 40:x + 76] + msg_key).digest()
3
4 aes_key = sha256_a[:8] + sha256_b[8:24] + sha256_a[24:32]
5 aes_iv = sha256_b[:8] + sha256_a[8:24] + sha256_b[24:32]
```

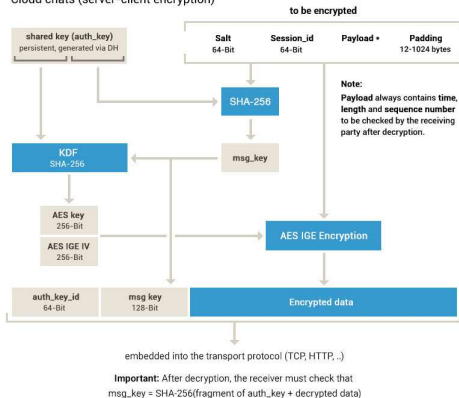
Auditing Crypto: MTPROTO 2

Usage of Weak Cryptographic Schemes

- KDF is not collision-resistant
- `auth_key[36: 40]` appear unused

MTPROTO 2.0, part I

Cloud chats (server-client encryption)



```
1 sha256_a = sha256(msg_key + auth_key[x: x + 36]).digest()
2 sha256_b = sha256(auth_key[x + 40:x + 76] + msg_key).digest()
3
4 aes_key = sha256_a[:8] + sha256_b[8:24] + sha256_a[24:32]
5 aes_iv = sha256_b[:8] + sha256_a[8:24] + sha256_b[24:32]
```

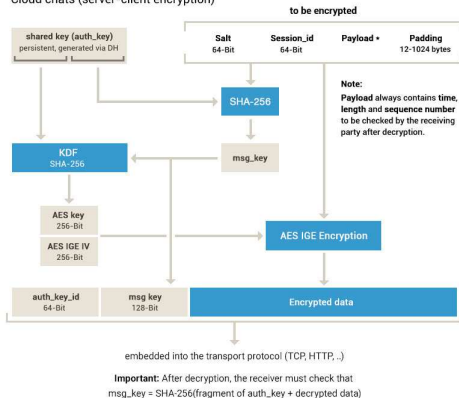

Auditing Crypto: MTPROTO 2

Replay attacks

- Auth key ID could be manipulated
- Old packet may be replayed

MTPROTO 2.0, part I

Cloud chats (server-client encryption)



Section 2

Auditing Implementations

Auditing Implementations: Pyrogram

- Use of naive RSA encryption
- Use of weak random generation with `random`
- Use of naive padding

Usage of Weak Cryptographic Schemes

- Use of naive RSA encryption

```
1 @classmethod
2 def encrypt(cls, data: bytes, fingerprint: int) -> bytes:
3     return int.to_bytes(
4         pow(
5             int.from_bytes(data, "big"),
6             cls.server_public_keys[fingerprint].e,
7             cls.server_public_keys[fingerprint].m
8         ),
9         256,
10        "big"
11    )
```

Usage of Weak Cryptographic Schemes

- Use of naive padding

```
1 padding = urandom(-(len(data) + len(sha)) % 255)
2 data_with_hash = sha + data + padding
3 encrypted_data = RSA.encrypt(data_with_hash, public_key_fingerprint)
```

Potential Side Channels

```
1 @classmethod
2 def ctr(cls, data: bytes, key: bytes, iv: bytearray, state: bytearray) -> bytes:
3     cipher = pyaes.AES(key)
4     ...
5
6     try:
7         iv[k] += 1
8         break
9     except ValueError:
10        iv[k] = 0
```

Section 3

A Custom Protocol

A Custom Protocol for Instant Messaging

- AE scheme
- Primitives
- Decide for low-level protocol(s) (TCP/UDP/WebRTC etc.)

- Level 1: Integrity and Privacy
- Level 2: + Packet Reordering
- Level 3: + Packet Replay
- Level 4: + Packet Loss

Setting

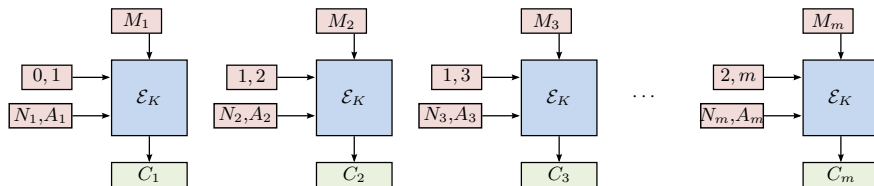
Environment:

- Probably: off-the-shelf 64-bit smartphones
- Messages: small messages
- Small associated data of up to 1kB
- Key(s) already exchanged

Security Goals:

Goal	Relevance	Level
Authentication of communication partners	essential	64-bit
Privacy of communication	essential	64-bit
Authenticity of communication	essential	64-bit
End-to-end encryption	essential	64-bit
Reorder, replay, and drop protection	essential	64-bit
Length hiding	essential	
Forward secrecy	essential	
Robustness against nonce misuse	optional	
Robustness against release of unverified plaintexts	optional	

Protocol Suggestion 1



- sAE Scheme: STREAM-like

- i = start/middle/end
- j = message index in a sequence

- AE Scheme: SIV-like

- misuse-resistant, messages are small and on-line not necessary

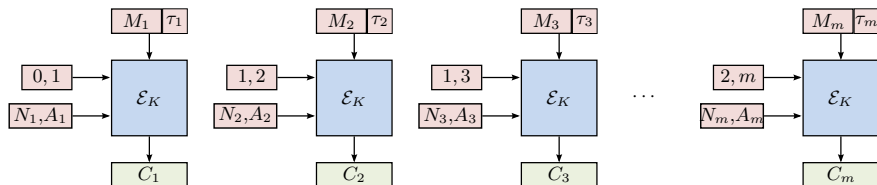
- Low-level protocols: TCP

- Ensures order and proper packet transport

```
1: function  $\mathcal{E}_K(i, j, N, A, M)$   
2:    $A' \leftarrow \text{ENCODEA}(i, j, A)$   
3:    $C \leftarrow \mathcal{E}_K^{N, A'}(M)$   
4:   return  $C$ 
```

```
1: function ENCODEA( $i, j, A$ )  
2:   return  $\langle i \rangle \parallel \langle j \rangle \parallel A$ 
```

Protocol Suggestion 2



- sAE Scheme: STREAM-like
 - τ_i = per-message stretch (length-hiding)
- AE Scheme: HCTR with AES-256 and PHash-AES-256
 - Tweakable STPRP with Encode-then-Encipher \implies Robust AE
 - Messages are small and on-line not necessary
 - Standardized single primitive
- Low-level protocols: TCP
 - Ensures order and proper packet transport

```
1: function  $\mathcal{E}_K(i, j, \tau, N, A, M)$   
2:    $V \leftarrow \text{ENCODEA}(i, j, \tau, N, A)$   
3:    $M' \leftarrow \text{ENCODEM}(M, \tau)$   
4:    $C \leftarrow \mathcal{E}_K^V(M')$   
5:   return  $C$ 
```

```
1: function  $\text{ENCODEA}(i, j, \tau, N, A)$   
2:   return  $\langle i \rangle \parallel \langle j \rangle \parallel \langle \tau \rangle \parallel N \parallel A$   
  
1: function  $\text{ENCODEM}(M, \tau)$   
2:   return  $M \parallel \langle \tau \rangle \parallel 0^\tau$ 
```

Questions?